

Quantum Information Theory

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Tue. H33 13pm c.t. & **Thu.** H34, 3pm c.t.
Mon. 12pm c.t., H33

Sheet 11

1 3-Qubit Code [10P]

Consider three qubits which are exposed to errors. The errors are rotations of the form $U = \exp(i\epsilon\sigma_x)$ with $\epsilon \ll 1$. Each qubit can be affected, thus

$$|\psi\rangle_E = E |\psi\rangle_L \quad \text{with} \quad E = U^{\otimes 3}. \quad (1)$$

The corresponding quantum circuit for encoding an arbitrary state $|\psi\rangle = \alpha|0\rangle + \beta|1\rangle$ with $\alpha, \beta \in \mathbb{C}$, $|\alpha|^2 + |\beta|^2 = 1$, and correcting the error is shown in Fig. 1.

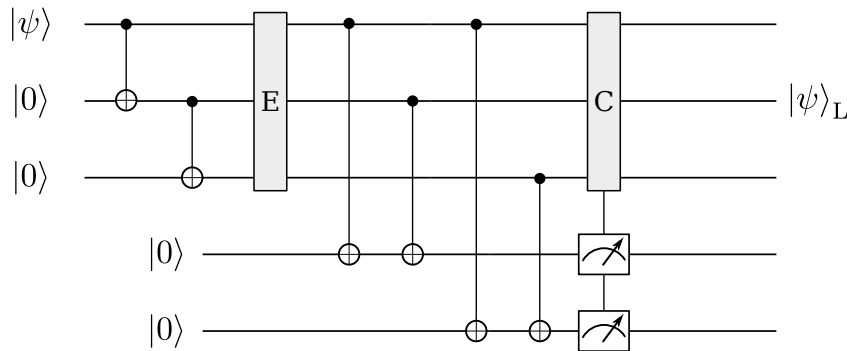


Figure 1: Quantum circuit for encoding and correcting (C-operation) a single bit-flip error (E-operation).

- Write down the ancilla measurements for a single σ_x error and the ambiguity of the syndrome results when multiple errors occur.
- Now, assume that the errors acting on the qubits are described by U . Corresponding to the measured ancilla block an appropriate correction operator is applied (see Fig. 1 and Tab. 1). Compare the fidelity $F_0 = |\langle\psi|U|\psi\rangle|^2$ without error correction code with the worst case fidelity (= state $|\psi\rangle_L$ is orthogonal to the state $\sigma_x^{\otimes 3} |\psi\rangle_L$) after applying the error correction.

ancilla	correction
00	$c_0 \psi\rangle_L + c_3 \sigma_x^{\otimes 3} \psi\rangle_L$
01	$c_1 \psi\rangle_L + c_2 \sigma_x^{\otimes 3} \psi\rangle_L$
10	$c_1 \psi\rangle_L + c_2 \sigma_x^{\otimes 3} \psi\rangle_L$
11	$c_1 \psi\rangle_L + c_2 \sigma_x^{\otimes 3} \psi\rangle_L$

Table 1: Quantum states after error correction. $c_0 = \cos^3(\epsilon)$, $c_1 = i \cos^2(\epsilon) \sin(\epsilon)$, $c_2 = -\cos(\epsilon) \sin^2(\epsilon)$, $c_3 = -i \sin^3(\epsilon)$.

2 Quantum operations [10P]

In the lecture the *operator-sum representation* has been used to represent quantum operations as $\sum_n E_n \rho E_n^\dagger$ using the *operation elements* $\{E_i\}$.

- (a) To get to this representation one assumes that the environment starts in a pure state. This can be done w.l.o.g. because we can *purify* a mixed state by introducing an additional system. To see this, consider two systems A and B with $\rho_A = \sum_n \lambda_n |a_n\rangle\langle a_n|$.
- (i) Show that $|\psi\rangle = \sum_n \sqrt{\lambda_n} |a_n\rangle_A \otimes |b_n\rangle_B$ is a *purification* of ρ_A for *any* orthonormal basis $\{|b_n\rangle\}_n$ with $\rho_A = \text{tr}_B(|\psi\rangle\langle\psi|)$.
- (ii) The purification of the system is not unique. Consider a second purifying system B' . Show that B and B' are related via a unitary transformation.
Hint: Write down the purified system using the Schmidt decomposition.
- (b) Consider the quantum circuit shown in Fig. 2. What are the corresponding operation elements E_n and what is the physical process described by this system?

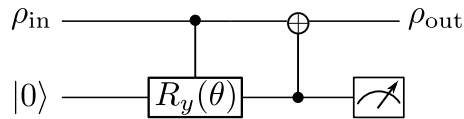


Figure 2: Quantum circuit

3 Generator Matrix and Hamming Code [8P]

Consider the generator matrix over a finite field of order 3¹

$$G = \begin{pmatrix} 0 & 1 & 2 & 1 \\ 1 & 0 & 1 & 0 \\ 1 & 2 & 2 & 1 \end{pmatrix}. \quad (2)$$

Calculate the corresponding *parity check matrix* H .

Hint: Transform G into the standard form, i.e., $G = [\mathbb{1} \ P]$.

¹also called Galois field, here $\text{GF}(3)$; the coset $\text{mod } 2$ is $\text{GF}(2)$