Universität Regensburg, Institut für Theoretische Physik Prof. Dr. Christoph Lehner, Andreas Hackl Winter 2022/23

# Exercise: Quantum Computing Problem set 2 (to be discussed in week of November 7, 2022)

#### Problem 1 Controlled-U

Show that

$$\begin{array}{c} \hline \\ A^{\dagger} \\ \hline \\ H \\ \end{array} \end{array} = \begin{array}{c} \hline \\ H \\ \hline \\ H \\ \end{array}$$

with

$$A = R_{\pi/2} H R_{\pi/4} \tag{1}$$

and



**Solution:** We first show that

$$AXA^{\dagger} = H \tag{2}$$

such that

$$A = \begin{pmatrix} \frac{1}{\sqrt{2}} & \frac{1+i}{2} \\ \frac{i}{\sqrt{2}} & \frac{1-i}{2} \end{pmatrix}$$
(3)

#### Problem 2 $C^n NOT$ with work qubits

The recursive definition of  $\mathbb{C}^n$ NOT given in the lecture has exponential cost for large n. If we have n-2 additional "work qubits", we can implement a gate whose cost only grows linearly with n. Show for n = 4 that



and then generalize this to n gates.

## Problem 3 Generalization of Deutsch-Jozsa algorithm

Let us consider a generalization of the Deutsch-Jozsa algorithm, where the input function f is not constrained to be either constant or balanced. Consider the scenarios of measuring r = 0 and  $r \neq 0$ . Show that measuring  $r \neq 0$  guarantees that the function is not constant and measuring r = 0 guarantees that the function is not balanced.

## Problem 4 Quantum parallelism

Write a circuit for  $U_f$  with N = 2 and  $f(x) = x \mod 2$ , i.e., f(x) = 0 if x is divisible by 2 and f(x) = 1 in all other cases. Generalize the circuit to general N.

## **Problem 5** Deutsch-Jozsa algorithm for N = 4 (optional)

Implement the Deutsch-Jozsa algorithm for the function of Problem 4 in the quantum computing simulator (http://github.com/lehner/sqc) for the case of N=4.